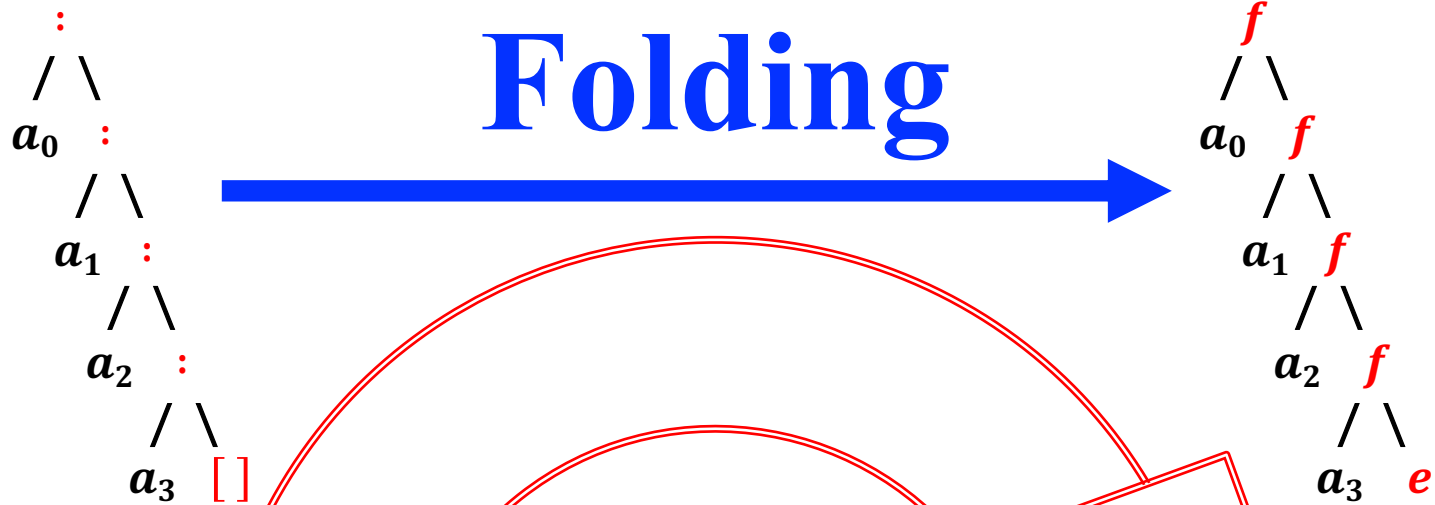


Folding



data *Nat* = *Zero* | *Succ Nat*

Common pattern for many recursive functions over *Nat* :

$f :: \text{Nat} \rightarrow \alpha$
 $f \text{ Zero} = c$
 $f (\text{Succ } n) = h (f n)$

$c :: \alpha$
 $h :: \alpha \rightarrow \alpha$

Three examples of such functions:

$(+) :: \text{Nat} \rightarrow \text{Nat} \rightarrow \text{Nat}$
 $m + \text{Zero} = m$
 $m + (\text{Succ } n) = \text{Succ } (m + n)$

$(\times) :: \text{Nat} \rightarrow \text{Nat} \rightarrow \text{Nat}$
 $m \times \text{Zero} = \text{Zero}$
 $m \times (\text{Succ } n) = (m \times n) + m$

$(\uparrow) :: \text{Nat} \rightarrow \text{Nat} \rightarrow \text{Nat}$
 $m \uparrow \text{Zero} = \text{Succ Zero}$
 $m \uparrow (\text{Succ } n) = (m \uparrow n) \times m$

The common pattern can be captured in a function:

$\text{foldn} :: (\alpha \rightarrow \alpha) \rightarrow \alpha \rightarrow \text{Nat} \rightarrow \alpha$
 $\text{foldn } h c \text{ Zero} = c$
 $\text{foldn } h c (\text{Succ } n) = h (\text{foldn } h c n)$

The three sample functions implemented using *foldn*:

$m + n = \text{foldn } \text{Succ } m n$
 $m \times n = \text{foldn } (\lambda x. x + m) \text{ Zero } n$
 $m \uparrow n = \text{foldn } (\lambda x. x \times m) (\text{Succ Zero}) n$

data *List* α = *Nil* | *Cons* α (*List* α)

Common pattern for many recursive functions over *List*:

$f :: \text{List } \alpha \rightarrow \beta$
 $f \text{ Nil} = c$
 $f (\text{Cons } x xs) = h x (f xs)$

$c :: \beta$
 $h :: \alpha \rightarrow \beta \rightarrow \beta$

Three examples of such functions:

$\text{sum} :: \text{List Nat} \rightarrow \text{Nat}$
 $\text{sum Nil} = \text{Zero}$
 $\text{sum } (\text{Cons } x xs) = x + (\text{sum } xs)$

$\text{length} :: \text{List } \alpha \rightarrow \text{Nat}$
 $\text{length Nil} = \text{Zero}$
 $\text{length } (\text{Cons } x xs) = \text{Succ Zero} + (\text{length } xs)$

$\text{append} :: \text{List } \alpha \rightarrow \text{List } \alpha \rightarrow \text{List } \alpha$
 $\text{append Nil } ys = ys$
 $\text{append } (\text{Cons } x xs) ys = \text{Cons } x (\text{append } xs ys)$

The common pattern can be captured in a function:

$\text{foldr} :: (\alpha \rightarrow \beta \rightarrow \beta) \rightarrow \beta \rightarrow [\alpha] \rightarrow \beta$
 $\text{foldr } f b \text{ Nil} = b$
 $\text{foldr } f b (\text{Cons } x xs) = f x (\text{foldr } f b xs)$

The three sample functions implemented using *foldr*:

$\text{sum } xs = \text{foldr } (+) \text{ Zero } xs$
 $\text{length } xs = \text{foldr } (\lambda x. \lambda n. n + (\text{Succ Zero})) \text{ Zero } xs$
 $\text{append } xs ys = \text{foldr } \text{Cons } ys xs$

Folding Unfolded

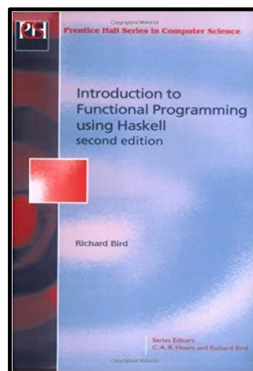
Polyglot FP for Fun and Profit
Haskell and Scala

See how **recursive functions** and **structural induction** relate to **recursive datatypes**

Follow along as the **fold abstraction** is introduced and explained

Watch as **folding** is used to simplify the definition of **recursive functions** over **recursive datatypes**

Part 1 - through the work of



Richard Bird

<http://www.cs.ox.ac.uk/people/richard.bird/>



Graham Hutton

[@haskellhutt](https://twitter.com/haskellhutt)

*A tutorial on the universality and
expressiveness of fold*

GRAHAM HUTTON

University of Nottingham, Nottingham, UK
<http://www.cs.nott.ac.uk/~gmh>



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