

Kleisli composition, flatMap, join, map, unit

a study/memory aid

to help learn/recall their implementation/interrelation

inspired by, and based on, the work of



Rob Norris

 @tpolecat



Bartosz Milewski

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<http://fpilluminated.com/>

This slide deck is meant both for (1) those who are familiar with the monadic functions that are **Kleisli composition**, **unit**, **map**, **join** and **flatMap**, and want to reinforce their knowledge (2) and as a memory aid, for those who sometimes need a reminder of how these functions are implemented and how they interrelate.

I learned about this subject mainly from **Functional Programming in Scala**, from **Bartosz Milewski's** YouTube videos (and his book), and from **Rob Norris's** YouTube video, **Functional Programming with Effects**.



Bartosz Milewski
@BartoszMilewski

Category Theory for Programmers



Bartosz Milewski



Rob Norris
@tpolecat



@philip_schwarz

See the final slide of this deck for some of the inspiration/ideas I got from **Rob Norris's** video.



If you need an intro to, or refresher on, the monadic functions that are **Kleisli composition**, **unit**, **map**, **join** and **flatMap**, then see the following



<https://www.slideshare.net/pjschwarz/rob-norrisfunctionalprogrammingwitheffects>

<https://www.slideshare.net/pjschwarz/kleisli-monad-as-functor-with-pair-of-natural-transformations>

<https://www.slideshare.net/pjschwarz/fish-operator-anatomy>

<https://www.slideshare.net/pjschwarz/kleisli-composition>

<https://www.slideshare.net/pjschwarz/compositionality-and-category-theory-a-montage-of-slidestranscript-for-sections-of-rnar-bjarnasons-keynote-composing-programs>

Runar Bjarnason

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**Functional Programming
In Scala**



Paul Chiusano

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A simple example of **hand-coding Kleisli composition** (i.e. `>=>`, the **fish operator**) for **Option** and **List**

Option

```
extension [A,B](f: A => Option[B])
  def >=>[C](g: B => Option[C]): A => Option[C] =
    a => f(a) match
      case Some(b) => g(b)
      case None => None
```

```
case class Insurance(name:String)
case class Car(insurance: Option[Insurance])
case class Person(car: Option[Car])
```

```
val car: Person => Option[Car] =
  person => person.car
```

```
val insurance: Car => Option[Insurance] =
  car => car.insurance
```

```
val carInsurance: Person => Option[Insurance] =
  car >=> insurance
```

```
val nonDriver= Person(car=None)
val uninsured = Person(Some(Car(insurance=None)))
val insured = Person(Some(Car(Some(Insurance("Acme")))))

assert(carInsurance(nonDriver).isEmpty)
assert(carInsurance(uninsured).isEmpty)
assert(carInsurance(insured).contains(Insurance("Acme")))
```

List

```
extension [A,B](f: A => List[B])
  def >=>[C](g: B => List[C]): A => List[C] =
    a => f(a).foldRight(List.empty[C]):
      (b, cs) => g(b) ++ cs
```

```
val toChars: String => List[Char] = _.toList
val toAscii: Char => List[Char] = _.toInt.toString.toList

assert( toChars("AB") == List('A','B') )
assert( toAscii('A') == List('6','5') )
```

```
val toCharsAscii: String => List[Char] =
  toChars >=> toAscii
```

```
assert(toCharsAscii("AB") == List('6','5','6','6'))
```



We have implemented `>=>` by hand. Twice. Once for **Option** and once for **List**.

Now let's make `>=>` generic and implement it in terms of the **built-in flatMap** function of the **Option** and **List Monads**.

Before

`>=>` is **hand-coded** and specialised for **Option** and **List**

```
extension [A,B](f: A => Option[B])
  def >=>[C](g: B => Option[C]): A => Option[C] =
    a => f(a) match
      case Some(b) => g(b)
      case None => None

extension [A,B](f: A => List[B])
  def >=>[C](g: B => List[C]): A => List[C] =
    a => f(a).foldRight(List.empty[C]):
      (b, cs) => g(b) ++ cs
```

After

`>=>` is **generic** and defined in terms of **built-in flatMap**

```
extension [A,B,F[_]: Monad](f: A => F[B])
  def >=>[C](g: B => F[C]): A => F[C] =
    a => f(a).flatMap(g)
```

```
trait Monad[F[_]]:
  def unit[A](a: => A): F[A]
  extension [A](fa: F[A])
    def flatMap[B](f: A => F[B]): F[B]

given Monad[Option] with
  def unit[A](a: => A): Option[A] = Some(a)
  extension [A](fa: Option[A])
    def flatMap[B](f: A => Option[B]): Option[B] = fa.flatMap(f)

given Monad[List] with
  def unit[A](a: => A): List[A] = List(a)
  extension [A](fa: List[A])
    def flatMap[B](f: A => List[B]): List[B] = fa.flatMap(f)
```



I first saw the definition of `>=>` in a YouTube video by **Bartosz Milewski**.



Bartosz Milewski's definition of the **fish operator** (**Kleisli composition**) in his lecture on **monads**.

$$f \gg= g = \lambda a \rightarrow \text{let } \underline{mb} = f\ a \text{ in } \underline{mb} \gg= g$$

See the next slide for more



 [@BartoszMilewski](#)  [Category Theory 10.1: Monads](#)



Kleisli Composition (fish operator)	<code>>=></code>	<code>compose</code>
Bind	<code>>>=</code>	<code>flatMap</code>
lifts <code>a</code> to <code>m a</code> (lifts <code>A</code> to <code>F[A]</code>)	<code>return</code>	<code>unit/pure</code>

class Monad m where

$(>=>) :: (a \rightarrow m b) \rightarrow (b \rightarrow m c) \rightarrow (a \rightarrow m c)$
return :: a -> m a

class Monad m where

$(>>=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$
return :: a -> m a



Kleisli Composition (fish operator) $>=>$ compose
Bind $>>=$ flatMap
lifts a to m a (lifts A to F[A]) return unit/pure

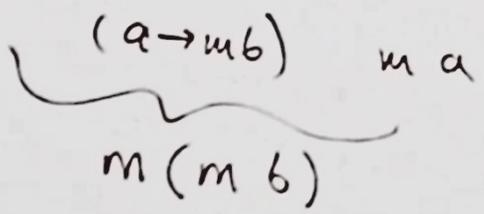
f g

$(>=>) :: (a \rightarrow m b) \rightarrow (b \rightarrow m c) \rightarrow (a \rightarrow m c)$

$f >=> g = \lambda a \rightarrow \text{let } \underline{mb} = f a \text{ in } \underline{mb} >>= g$

$(>>=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$

$ma >>= f = \text{join } (fmap f ma)$



join :: m (m a) -> m a

class Monad m where

$(>>=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$

return :: a -> m a

class Functor m => Monad m where

join :: m(m a) -> m a

return :: a -> m a

$>=>$ Kleisli Composition (aka the fish operator)

$>>=$ Bind

$f >=> g = \lambda a \rightarrow \text{let } mb = f a \text{ in } mb >>= g$
 $= \lambda a \rightarrow (f a) >>= g$

class Functor f where

fmap :: (a -> b) -> f a -> f b

class Functor m => Monad m where

join :: m(m a) -> m a

return :: a -> m a



Now let's hand-code ourselves the **flatMap** functions of the **Option** and **List Monads**.

```
trait Monad[F[_]]:  
  def unit[A](a: => A): F[A]  
  extension [A](fa: F[A])  
    def flatMap[B](f: A => F[B]): F[B]
```

```
extension [A,B,F[_]: Monad](f: A => F[B])  
  def >=>[C](g: B => F[C]): A => F[C] =  
    a => f(a).flatMap(g)
```

Before

>=> is **generic** and defined in terms of **built-in flatMap**

```
given Monad[Option] with  
  def unit[A](a: => A): Option[A] = Some(a)  
  extension [A](fa: Option[A])  
    def flatMap[B](f: A => Option[B]): Option[B] = fa.flatMap(f)
```

```
given Monad[List] with  
  def unit[A](a: => A): List[A] = List(a)  
  extension [A](fa: List[A])  
    def flatMap[B](f: A => List[B]): List[B] = fa.flatMap(f)
```

After

>=> is **generic** and defined in terms of **hand-coded flatMap**

```
given Monad[Option] with  
  def unit[A](a: => A): Option[A] = Some(a)  
  extension [A](fa: Option[A])  
    def flatMap[B](f: A => Option[B]): Option[B] =  
      fa match  
        case Some(a) => f(a)  
        case None => None
```

```
given Monad[List] with  
  def unit[A](a: => A): List[A] = List(a)  
  extension [A](fa: List[A])  
    def flatMap[B](f: A => List[B]): List[B] =  
      fa.foldRight(List.empty[B]):  
        (a,bs) => f(a) ++ bs
```



 @philip_schwarz

Earlier we implemented a generic `>=>` in terms of the **built-in flatMap** function of the **Option** and **List Monads**.

Let's do that again but this time implementing `>=>` in terms of the built-in **map** and **join** functions of the **Option** and **List Monads**.

The next slide is just a refresher on the fact that it is possible to define a **Monad** in terms of **unit**, **map** and **join**, instead of in terms of **unit** and **flatMap**, or in terms of **unit** and the **fish operator**. If it is the first time that you go through this slide deck, you may want to skip the slide.



in Scala, **join** is called **flatten**

Bartosz Milewski introduces a third definition of Monad in terms of join and return, based on Functor

The whiteboard contains the following handwritten text:

$(\Rightarrow) :: (a \rightarrow m b) \rightarrow (b \rightarrow m c) \rightarrow (a \rightarrow m c)$ #1

$f \Rightarrow g = \lambda a \rightarrow \text{let } \underline{mb} = f a \text{ in } \underline{mb} \gg= g$

$(\gg=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$

$ma \gg= f = \text{join } (fmap f ma)$

$\text{join} :: m (m a) \rightarrow m a$

class Monad m where

$(\gg=) :: m a \rightarrow (a \rightarrow m b) \rightarrow m b$ #2

$\text{return} :: a \rightarrow m a$

class Functor m => Monad m where

$\text{join} :: m (m a) \rightarrow m a$ #3

$\text{return} :: a \rightarrow m a$

So this (**join** and **return**) is an alternative definition of a monad. But in this case I have to specifically say that m is a **Functor**, which is actually a nice thing, that I have to explicitly specify it.

...

But remember, in this case (**join** and **return**) you really have to assume that it is a **functor**. In this way, join is the most basic thing. Using just join and return is really more atomic than using either bind or the Kleisli arrow, because they additionally **subsume functoriality**, whereas here, functoriality is separate, separately it is a functor and separately we define **join**, and separately we define **return**.

...

So this definition (**join** and **return**) or the definition with the **Kleisli arrow**, they are not used in **Haskell**, although they could have been. But **Haskell** people decided to use this (**>>=** and **return**) as their basic definition and then for every monad they separately define **join** and the **Kleisli arrow**. So if you have a monad you can use **join** and the **Kleisli arrow** because they are defined in the library for you. So it's always enough to define just **bind**, and then **fish** and **join** will be automatically defined for you, you don't have to do it.

Before

`>=>` is **generic** and defined in terms of **built-in flatMap**

```
trait Monad[F[_]]:  
  def unit[A](a: => A): F[A]  
  extension [A](fa: F[A])  
    def flatMap[B](f: A => F[B]): F[B]  
  
extension [A,B,F[_]: Monad](f: A => F[B])  
  def >=>[C](g: B => F[C]): A => F[C] =  
    a => f(a).flatMap(g)
```

```
given Monad[Option] with  
  def unit[A](a: => A): Option[A] = Some(a)  
  extension [A](fa: Option[A])  
    def flatMap[B](f: A => Option[B]): Option[B] = fa.flatMap(f)
```

```
given Monad[List] with  
  def unit[A](a: => A): List[A] = List(a)  
  extension [A](fa: List[A])  
    def flatMap[B](f: A => List[B]): List[B] = fa.flatMap(f)
```

After

`>=>` is **generic** and defined in terms of **built-in map** and **join**

```
trait Functor[F[_]]:  
  extension [A](fa: F[A])  
    def map[B](f: A => B): F[B]  
  
trait Monad[F[_]] extends Functor[F]:  
  def unit[A](a: => A): F[A]  
  extension [A](ffa: F[F[A]])  
    def join: F[A]  
  
extension [A,B,F[_]: Monad](f: A => F[B])  
  def >=>[C](g: B => F[C]): A => F[C] =  
    a => f(a).map(g).join
```

```
given Monad[Option] with  
  def unit[A](a: => A): Option[A] = Some(a)  
  extension [A](fa: Option[A])  
    def map[B](f: A => B): Option[B] = fa.map(f)  
  extension [A](ffa: Option[Option[A]])  
    def join: Option[A] = ffa.flatten
```

```
given Monad[List] with  
  def unit[A](a: => A): List[A] = List(a)  
  extension [A](fa: List[A])  
    def map[B](f: A => B): List[B] = fa.map(f)  
  extension [A](ffa: List[List[A]])  
    def join: List[A] = ffa.flatten
```



Now let's hand-code ourselves the **map** and **join** functions of the **Option** and **List Monads**.

```
trait Functor[F[_]]:  
  extension [A](fa: F[A])  
    def map[B](f: A => B): F[B]
```

```
trait Monad[F[_]] extends Functor[F]:  
  def unit[A](a: => A): F[A]  
  extension [A](ffa: F[F[A]])  
    def join: F[A]
```

```
extension [A,B,F[_]: Monad](f: A => F[B])  
  def >=>[C](g: B => F[C]): A => F[C] =  
    a => f(a).map(g).join
```

Before

>=> is **generic** and defined in terms of **built-in map** and **join**

```
given Monad[Option] with  
  def unit[A](a: => A): Option[A] = Some(a)  
  extension [A](fa: Option[A])  
    def map[B](f: A => B): Option[B] = fa.map(f)  
  extension [A](ffa: Option[Option[A]])  
    def join: Option[A] = ffa.flatten
```

```
given Monad[List] with  
  def unit[A](a: => A): List[A] = List(a)  
  extension [A](fa: List[A])  
    def map[B](f: A => B): List[B] = fa.map(f)  
  extension [A](ffa: List[List[A]])  
    def join: List[A] = ffa.flatten
```

After

>=> is **generic** and defined in terms of **hand-coded map** and **join**

```
given Monad[Option] with  
  def unit[A](a: => A): Option[A] = Some(a)  
  extension [A](fa: Option[A])  
    def map[B](f: A => B): Option[B] = fa match  
      case Some(a) => Some(f(a))  
      case None => None  
  extension [A](ffa: Option[Option[A]])  
    def join: Option[A] = ffa match  
      case Some(a) => a  
      case None => None
```

```
given Monad[List] with  
  def unit[A](a: => A): List[A] = List(a)  
  extension [A](fa: List[A])  
    def map[B](f: A => B): List[B] =  
      fa.foldRight(List.empty[B])((a,bs) => f(a)::bs)  
  extension [A](ffa: List[List[A]])  
    def join: List[A] =  
      ffa.foldRight(List.empty[A])((a,as) => a ++ as)
```



I would like to thank **Rob Norris** for his great talk (see next slide), from which I learned a lot about **Kleisli composition** and in which I first saw the use of a **syntax class** to add the **fish operator** to a type class.



scale.bythebay.io

Rob Norris

Functional Programming with Effects

```
// A typeclass that describes type constructors that allow composition with >=>
trait Fishy[F[_]] {

  // Our identity, A => F[A] for any type A
  def pure[A](a: A): F[A]

  // Composition - the "fish" operator
  def >=>[A, B, C](f: A => F[B], g: B => F[C]): A => F[C] =
    a => f(a).flatMap(g) // hey that looks like flatMap!
}
```

the **fish operator** (Kleisli composition), can be implemented using **flatMap**.



```
// A typeclass that describes type constructors that allow composition with >=>
trait Fishy[F[_]] {

  // Our identity, A => F[A] for any type A
  def pure[A](a: A): F[A]

  // The operation we need if we want to define >=>
  def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

We can implement **compose**, the **fish operator** using **flatMap**, so the **fish operator** is something we can derive later really, the operation we need is **flatMap**.

```
// Now we can define >=> as an infix operator using a syntax class
implicit class FishyFunctionOps[F[_], A, B](f: A => F[B]) {
  def >=>[C](g: B => F[C])(implicit ev: Fishy[F]): A => F[C] =
    a => ev.flatMap(f(a))(g)
}

// Let's define an instance for Option
implicit val FishyOption: Fishy[Option] =
  new Fishy[Option] {
    def pure[A](a: A) = Some(a)
    def flatMap[A, B](fa: Option[A])(f: A => Option[B]) = fa.flatMap(f)
  }
```

We can define a **syntax class** that adds these methods so that anything that is an F[A], if there is a **Fishy** instance, gets these operations by syntax.



And this **Fishy** typeclass that we have derived from nothing, using math, is **Monad**. So this scary thing, it just comes naturally and I haven't seen people talk about getting to it from this direction. And so I hope that was helpful.

Fishy

```
// Fishy typeclass
trait Fishy[F[_]] {
  def pure[A](a: A): F[A]
  def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```

Monad

```
// Monad typeclass
trait Monad[F[_]] {
  def pure[A](a: A): F[A]
  def flatMap[A, B](fa: F[A])(f: A => F[B]): F[B]
}
```